Du Bois Reymond à la Conway (the game)

BY VINCENT BAGAYOKO (IMJ-PRG, PARIS)

Colloquium Logicum, ÖAW Vienna, 10-09-24





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increasing, have Taylor expansions and inverses	

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$$\partial : \mathbf{No} \longrightarrow \mathbf{No}$$

and a composition law

$$\circ: \mathbf{No} \times \mathbf{No}^{>\mathbb{R}} \longrightarrow \mathbf{No}$$

such that for each $a \in \mathbf{No}$, the function $\hat{a} : \mathbf{No}^{>\mathbb{R}} \longrightarrow \mathbf{No}$ behaves like a germ in a Hardy field.

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For instance, each \hat{a} should be strictly monotonous and differentiable with $\hat{a}' = \hat{\partial}(a)$, we should have Taylor expansions, and for fixed $\xi \in \mathbf{No}^{>\mathbb{R}}$, the function $a \mapsto a \circ \xi$ should be an endomorphism of $(\mathbf{No}, +, \cdot, <)$.

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Let's play a game instead.

Rules of the game:

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 My goal is to convince you that I know what the function â : No^{>ℝ} → No; ξ → a ∘ ξ should be.

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If you win, I'll buy you a drink before the end of times.

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Reals and small ordinals

The ordered field of surreal numbers contains a canonical copy of the ordered field of real numbers:

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 ${\bf No}$ also contains the ordered semi-ring ${\bf On}$ of ordinal numbers under the commutative natural/Hessenberg arithmetic.

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For instance $\omega = \{\mathbb{N} \mid \emptyset\}$ is the simplest positive infinite number. It should correspond to the simplest germ that tends to $+\infty$, i.e. to the identity function $\hat{\omega} = \xi \mapsto \xi$.

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So elements of $\mathbb{R}(\omega)$ should act as the corresponding rational functions. What else?

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Suppose that α is the β -th such fixed point. We then⁹ define α using $\hat{\beta}$ and the ω^{η} -th iterate $\exp_{\omega^{\eta}}$ of the exponential...

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If $(\mathfrak{m}_{\gamma})_{\gamma < \eta}$ is a strictly decreasing sequence of monomials and $(r_{\gamma})_{\gamma < \eta} \in \mathbb{R}^{\eta}$, then one defines inductively the sum

$$\sum_{\gamma < \eta} r_{\gamma} \mathfrak{m}_{\gamma} = \left\{ \sum_{\gamma < \rho} r_{\gamma} \mathfrak{m}_{\gamma} + q \mathfrak{m}_{\rho} : q \in (-\infty, r_{\rho}) \mid \sum_{\gamma < \rho} r_{\gamma} \mathfrak{m}_{\gamma} + q \mathfrak{m}_{\rho} : q \in (r_{\rho}, +\infty) \right\}.$$

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Claim: I can define each $\widehat{\mathfrak{m}}_{\gamma}$ in so that for all $\xi \in \mathbf{No}^{>\mathbb{R}}$, the following number is well-defined:

$$\hat{a}(\xi) := \mathfrak{n} \longmapsto \sum_{\gamma < \eta} r_{\gamma} \, \widehat{\mathfrak{m}}_{\gamma}(\xi)(\mathfrak{n}).$$

Have you stumbled upon $\omega^{\frac{1}{\omega}} = \left\{ \mathbb{N} \mid \left\{ \omega^{\frac{1}{n}} : n \in \mathbb{N}^{>0} \right\} \right\}$? You must have a retorse mind.

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This function was defined by Gonshor ('86).

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... but also has Taylor expansions and behaves like a germ in a Hardy field? Hardy-type asymptotics entail that the simplest choice for $\widehat{\omega}^{\omega}(\xi + \zeta)$ is $\widehat{\omega}^{\omega}(\xi) \widehat{\omega}^{\omega}(\zeta)$. Gonshor ('86) defined the exponential function as follows, for $\xi = \{L \mid R\}$

$$\exp(\xi) = \left\{ \exp(l) \sum_{i \leqslant n} \frac{(\xi - l)^i}{i!} \mid \frac{\exp(r)}{i!} \right\}$$

where l, r, i range in L, R and \mathbb{N} respectively.

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If one is lucky, the number φ is "simpler" in some sense than \mathfrak{m} , so I can claim inductively that I know what $\hat{\varphi}$ is. And then $\hat{\mathfrak{m}} = \exp \circ \hat{\varphi}$. With a bit less luck, going through monomials in φ and the inductively so, we end up finding only simpler numbers after some time.

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But that may not happen. One may not reach any simpler number.

Theorem. [BERARDUCCI-MANTOVA, 2017] If the process continues indefinitely, then at some stage, problematic monomials \mathfrak{n} thus appearing have the form

$$\mathfrak{n} = \mathrm{e}^{\varphi_1 \pm \mathrm{e}^{\varphi_2 \pm \mathrm{e}^{\varphi_3 \pm \mathrm{e}^{\cdot}}}}$$

for some strictly simpler $\varphi_1, \ldots, \varphi_n, \ldots$

You seem to have found a nested number. Well done! those are the worst.

Each monomial $\mathfrak{m} \in \mathbf{No}$ (do you know about the Hahn series deccomposition yet?) can be written as $\mathfrak{m} = e^{\varphi}$ for $\varphi = \log \mathfrak{m}$ a surreal number of a special type.

If one is lucky, the number φ is "simpler" in some sense than \mathfrak{m} , so I can claim inductively that I know what $\hat{\varphi}$ is. And then $\hat{\mathfrak{m}} = \exp \circ \hat{\varphi}$. With a bit less luck, going through monomials in φ and the inductively so, we end up finding only simpler numbers after some time.

But that may not happen. One may not reach any simpler number.

Theorem. [BERARDUCCI-MANTOVA, 2017] If the process continues indefinitely, then at some stage, problematic monomials \mathfrak{n} thus appearing have the form

$$\mathfrak{n} = \mathrm{e}^{\varphi_1 \pm \mathrm{e}^{\varphi_2 \pm \mathrm{e}^{\varphi_3 \pm \mathrm{e}^{\cdot}}}}$$

for some strictly simpler $\varphi_1, \ldots, \varphi_n, \ldots$.

With van der Hoeven, we extended this to all numbers using hyperexponentials and their inverses.

So how do I deal with nested monomials? "An" example is

$$\mathfrak{n} = e^{\sqrt{\omega} + e^{\sqrt{\log \omega}} + e^{\sqrt{\log \log \omega}} + e^{\cdot}}$$

Suffices to deal with
$$n_i = e^{\sqrt{\log_i \omega} + e^{\sqrt{\log_{i+1} \omega} + e^{-\frac{1}{2}}}}$$
 for some $i \in \mathbb{N}$.

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For $\xi \in \mathbf{No}^{>\mathbb{R}}$, what should $\tilde{\mathfrak{n}}_i(\xi)$ be? The answer comes from a simplicity heuristic (a choice) and Taylor expansions. Rough idea:

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- Choosing *i* large enough, one can insure that φ is longest, such that the Taylor expansion of $\hat{\mathfrak{n}}_i(\varphi)$ with radius ε converges formally. We thus set $\hat{\mathfrak{n}}_i(\xi) := \sum_{k \in \mathbb{N}} \frac{\hat{\mathfrak{n}}_i^{(k)}(\varphi)}{k!} \varepsilon^k$.

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I claim that $\hat{\varepsilon_0}$ is a surreal version of ABEL/KNESER/ECALLE's ω -th iterate E of exp, a realanalytic function on $\mathbb{R}^{\geq 0}$ satisfying $E(t+1) = \exp(E(t))$ for all $t \geq 0$.

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Indeed since $\varepsilon_0 = \{\{\omega, \exp(\omega), \exp(\exp(\omega)), \dots\} \mid \emptyset\}$, the function $\hat{\varepsilon_0}$ should grow faster than all finite iterates of exp. Hardy-type asymptotics give an approximation of

 $\hat{\varepsilon}_0' \approx \hat{\varepsilon}_0 (\log \circ \hat{\varepsilon}_0) (\log \circ \log \circ \hat{\varepsilon}_0) \cdots = T \circ \hat{\varepsilon}_0$ for a transseries T.

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So if we can define $\hat{\varepsilon_0}$ at sufficiently simple φ 's, then we can extend the definition to all $\varphi + \varepsilon$ for sufficiently small ε , by

$$\hat{\varepsilon}_0(\varphi + \varepsilon) := \sum_{k \in \mathbb{N}} \frac{\hat{\varepsilon}_0^{(k)}(\varphi)}{k!} \varepsilon^k.$$

Some heuristics suggest that the simplest value for $\hat{\varepsilon}_0(\omega+1)$ is $\exp(\hat{\varepsilon}_0(\omega))$. Thus we should have $\hat{\varepsilon}_0(\xi+1) = \exp(\hat{\varepsilon}_0(\xi))$ for all $\xi \in \mathbf{No}^{>\mathbb{R}}$.

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To define $\hat{\varepsilon}_0$ at a ξ , it suffices to define it at $\xi - n$ for some $n \in \mathbb{N}$ and then take the *n*-fold iterated exponential of the result. Fix $\xi \in \mathbf{No}^{>\mathbb{R}}$ and suppose $\hat{\varepsilon}_0(\zeta)$ is defined for simpler ζ .

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If that $\xi - n$ has a strict (hence simpler) truncation φ with $\xi - n - \varphi < \frac{\hat{\varepsilon}_0(\varphi)}{\hat{\varepsilon}_0'(\varphi)}$ for some $n \in \mathbb{N}$, then we define $\hat{\varepsilon}_0(\xi)$ via Taylor expansions.

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So it suffices to define $\hat{\varepsilon_0}$ on the class \mathbf{Tr} of numbers that do not have such truncations. For $\varphi, \psi \in \mathbf{Tr}, \varphi < \psi$, monotonicity of $\hat{\varepsilon_0}$ entails that $\mathcal{E}_n^{\pm}(\hat{\varepsilon_0}(\varphi)) < \hat{\varepsilon_0}(\psi)$ where

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This gives an inductive definition, if φ is the simplest element of **Tr** with $L < \varphi < R$, $L, R \subseteq$ **Tr**:

$$\hat{\varepsilon_0}(\varphi) = \{ \exp^{\circ n}(\varphi), \mathcal{E}_n^+(\hat{\varepsilon_0}(l)) : n \in \mathbb{N} \land l \in L \} \mid \{ \mathcal{E}_n^-(\hat{\varepsilon_0}(r)) : n \in \mathbb{N} \land r \in R \}.$$

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General case of $\exp_{\omega^{\mu}}$: proceed inductively, replacing $\{\exp^{\circ n} \in \mathbb{N}\}$ by $\{\exp^{\circ n}_{\omega^{\eta}}: n \in \mathbb{N} \land \eta < \mu\}$.

I win the game. Thanks for playing along!



Proof of VAN DER HOEVEN'S CONJECTURE. Proof by game. If the conjecture were false, then 10 minutes should be enough, to a room full of smart people, to disprove it. But they just lost the very fair game. Therefore the conjecture is true.

Congrats I guess

You win!



Thanks for playing along (and ruining my project; saves time)!